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#### INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(51) International Patent Classification <sup>5</sup>:

H04N 7/173

(11) International Publication Number:

WO 92/15174

7/173

(43) International Publication Date:

3 September 1992 (03.09.92)

(21) International Application Number:

PCT/GB92/00337

A1

(22) International Filing Date:

25 February 1992 (25.02.92)

(30) Priority data:

9103907.3

25 February 1991 (25.02.91) GB

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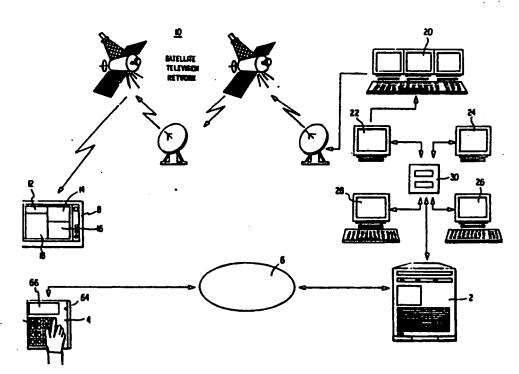
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(81) Designated States: AT (European patent), BE (European patent), CH (European patent), DE (European patent), DK (European patent), ES (European patent), FR (European patent), GB (European patent), GR (European patent), IT (European patent), JP, LU (European patent), MC (European patent), NL (European patent), SE (European patent).

#### **Published**

With international search report.

(54) Title: INTERACTIVE TRANSACTION PROCESSING SYSTEM



(57) Abstract

A distributed bject-based transaction system provides a plurality of terminals (4) and/or host computers (2) n which objects, or named memory spaces, reside. The bjects are controlled by methods which are located on each respective terminal r host on which an bject resides, and the methods can be invoked by any of the terminals or host computers. Updating the objects is accomplished by invoking the relevant method on each of the nodes wherein the object resides. The distributed object-based transaction system is particularly useful in the implementation f an auction system for remotely situated bidders utilizing interactive televisi n.

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### INTERACTIVE TRANSACTION PROCESSING SYSTEM

#### TECHNICAL FIELD

This invention relates to the art of data processing and th combination of data processing with video displays of items relat d to the data. In the ultimate preferred embodiment, the invention is an auction system for remotely situated bidders conduct d utilizing interactive television.

#### BACKGROUND

The invention is an interactive transaction processing syst m and data processing method. The fundamental system is a distributed object-based transaction system having a wide rang of potential uses. This distributed object-based transaction system has particular utility in the implementation of an interactive televised auction in which remote subscribers bid in competition with the live auction in the saleroom. Thus, disclosure of the invention will be facilitated by first providing a basic appreciation of the operation of an auction.

An auction normally proceeds under the control of an auctioneer. An item to be auctioned (sold) is displayed, and the auctioneer asks for a "bid" for the item. The auctioneer can set an opening bid price. The auctioneer invites further bids without

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specifying the next pric. The auctioneer states the pric who is held up to signal to the auctioneer that the bidder wishes to bid. The auctioneer usually accepts the bid of the first bidder to hold up his paddle and states the new price. The number on the paddle of the bidder whose bid has been accepted is recorded by the auctioneer or his clerk.

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Saleroom etiquette dictates that if the auctioneer senses that two bidders are competing for the item being sold, he should conduct a "ping-pong" auction between these two bidders. A ping-pong auction is a process whereby the auctioneer ignores bids by bidders other than the chosen two bidders competing for the it m. When one of the ping-pong bidders drops out, e.g., by failing to make a bid, the auctioneer will then accept bids from oth r bidders.

The system of the invention is designed to allow an auction in accordance with this process to be conducted at a plurality of remote sites. In general, terms used in the art of auctions will be used to describe the invention, and the bidders at the rem t sites will be referred to as "subscribers". It will b appreciated, however, that the invention is equally applicable to a variety of operations other than auctions.

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#### SUMMARY OF THE INVENTION

### A. Distributed Object-based Data Processing

In acc rdance with the inv nti n, a distributed obj ct-bas d data processing program provides a system for handling a br ad range of transactions and is specifically adapted to conduct an auction in the preferred embodiment. In accordance with th technique known as object-oriented design, the distributed object-based data processing program comprises "methods", "objects", and "classes".

An "object" is a data element, the specific definition of which depends upon the desired transaction or other operation. In the auction system which will be described in detail below, an example of an object is the "bid display." This is a list of identifications of the first ten bidders whose bids were received by the host computer in order of receipt of the bids.

A "method" is a set of instructions which are provided to the host computer to tell the computer what processes are to be performed on the object. For example, when a bid is received, the method "update bid display" is invoked which adds the identification of a bidder (another object) to the bid display or cancels a bid.

A "class" is a group of related methods, or routines. For example, the "user display class" includes the "add user" and "update user" methods. Classes can be arranged in a structure where one class "inherits" methods from another class (termed its

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"superclass"). The ultimate "ancestor" class is the "root" class, which contains the most generally applicable methods. The methods are grouped into classes to maximiz efficiency, minimiz the impact of subsequent design changes and promote reusability of code. Also, the bulk of the source code is reduced since terms common to the grouped methods do not have to be redefined for each of the related methods.

The typical instruction to the computer, which can be programmed in any of a variety of languages is: "Perform method Be on object A". This would cause the computer first to find the class of object A by searching a table of classes. Then,—it searches the method table of that particular class to confirm that method B is indeed a part of that class. If the computer cann the find the particular method in the given class, it searches its superclass. The computer then invokes the particular method by calling a subroutine identified by the method and performing the steps which comprise the method.

After the particular method has been performed on the object, the system updates the value of that object at all nodes in which the particular object resides by invoking the same method at this nodes. Thus, if the object "bid display" resides on seviral workstations and the host, the new value of the "bid display" object resulting from the completion of an invoked method relating to that object will be propagated to all of the nodes by the updat

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proc ss which invokes the sam method at each f th remaining nod s having th "bid display" obj ct.

Th distributed object-bas d transaction system of invention integrates an application which is divided into several processes and running on several machines. For example, an application may be terminal-based and intended for implementati n on a single host computer, or it may be implemented on multiple workstations connected to a variety of other devices. auction example described below, personal computers have programs capable of implementing individual methods relating to the obj cts located on the particular personal computer. The invoking of a method, which inherently changes the value of an object, is always automatically communicated to the other workstations and the host computer having that object by the process of updating whereby that method in also invoked at all nodes having that object. Thus, th distributed object-based transaction system provides a uniform mechanism for interaction between application components whil allowing each platform to contribute maximum functionality.

Remote procedure call mechanisms are known and give the programmer the ability to call a subroutine in the normal way but have it execute in another process that could be located in anoth r machine. To the programmer, the result appears in exactly the sam as if the subroutine had been part of his program and had been executed in the same process.

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The remote procedur call mechanism used in the system of the invention, as well as that of other systems, such as Sun's RPC, employs a signal description of the real routine in the caller's code. The stub routine then communicates with the remote process and passes the necessary information to it so that the correct routine is invoked in the remote process. The remote process will then return any results of the subroutine call back to the stub routine, which will be waiting for the reply. The stub routine then unpacks the return definition and places it into the proper variables. The stub routine then returns control to the calling routine and execution continues as normal.

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The distributed object-based transaction system of th invention implements a remote procedure call mechanism and rout s requests and responses across external networks and internal int r-process communications structures (such as pipes and queues). The system can cope with multiple paths between two nodes and chooses the most efficient route available.

The distributed object-based transaction system allows th combination of an advanced, general purpose application architecture with specific support for host computer features.

By permitting any node (object location) to be either a cli nt or a server, the distributed object-based transaction system transcends fixed client-server roles for networked computers. In the traditional system, a server offers a range of services and the

client access s those s rvices. In the distributed object-based transaction syst me of the invention, the server has a range of bj cts, and the client is on wheap accesses those objects. Any of the workstations may be a client with respect to some objects as well as a server with respect to other objects. This allows workstations, for instance, to be informed of dynamically changing information as well to initiate their own transactions.

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The distributed object-based transaction system disclosed herein has the further advantage that all interactions between nodes on the network are based on a single model. The mod 1 follows the object-oriented paradigm and has the advantage of local transparency to the clients in that the client need only identify the data object in a call and not the servers associated with that object. For example, a read only request will be directed only t the server nearest the client, whereas an update would be propagated to all locations of the object.

A basic feature of the distributed object-based transaction system of the invention is that the data can be replicated becaus the objects can reside in multiple nodes. In the traditional system, the host is generally the server because it contains the database. In contrast, the distributed object-based transaction system allows each of the workstations to have a database related to the objects residing on that workstation. The system ensures consistency between multiple copies of the same data by routing update requests to all relevant servers for a given object. For

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example, if a number of workstati ns display the valu of a particular data bj ct which als resides on the host, the system will invoke methods in each of the involved workstations t cause all of the values to be the same in all workstations concern d. There is, thus, no requirement to explicitly code for th distribution of the data because the system automatically updates the value of the object at all nodes wherein that object is located.

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Transaction management, like shared-memory and concurr nt processing support, is a specific feature of the Stratus machine, which is the preferred host computer. While equivalent facilities exist in a few other environments the distributed object-bas d system of the invention manages the Stratus version of them. Transaction protection ensures that every transaction will eith r complete successfully or be fully "backed-out", i.e., leave n trace that it ever executed. This feature is vital for transaction procession systems and is extremely complex to emulate if n t available.

Without transaction protection, a database can become inconsistent if 1) two transactions interfere with each other, e.g. by updating the same record or 2) a transaction fails during execution leaving some updates performed but not others. A transaction protection system will ensure that two transactions d not interfere with each other and that, if one does not complete, it is fully backed-out.

The pr grammer normally has to call subroutines to indicate the start and nd f transactions. Ending a transacti n n rmally is call d 'committing' because at this point th changes made tak effect apparently instantaneously and cannot thereafter be undon. In the transaction, the programmer must check the status of every file operation to determine whether the transaction can proceed or whether the transaction protection system has detected a conflict. If a problem is detected, the programmer must abort the transaction, i.e. end it abnormally. Ideally, because conflicts can arise in the normal course of events, the programmer should attempt to execute the entire transaction again because the cause of the conflict (another transaction) will finish eventually.

Instead of the programmer having to code for these functions, the distributed object-based transaction system of the invention takes over the management of transactions. Because a transacti n in accordance with the system of the invention corresponds to a method, the system will start the transaction before calling th method code and, when the method finishes, the system can commit. The system also has all the information at hand to restart transactions which fail. It does this by offering the programm r equivalents to all the file operation subroutines. These equivalents call the real file subroutine but check the status of the result. If they detect a conflict, the method is aborted at that point and control jumps back to the system which can call the transaction again automatically.

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An ther feature f the system of the invention is its port p 1 management. Before the data of a fill can be accessed in a program, a port must be "open d" to that file. There is an "pen" operating system routine which locates the file by name and creat s a port through which that program can access it. Thereafter, the file is accessed by the given port number, not by its name. In the open call, the programmer specifies what kind of access is required, such as input or update, indexed or sequential. The port also remembers the program's current position in the file, so that calls can be made to read successive records or update the current record.

A batch program will open ports, perform file operations and then close them again. With an on-line, real-time system, p rts cannot be opened for each transaction as this would create an unacceptable overhead. Instead, ports are opened once when th system starts and stay open while the system is up.

Because the system of the invention allows methods to be called from anywhere, including from other methods, a situation can arise where a method uses the same port as the method calling it. If, in using the port, the method altered the current position in the file then the calling method would lose its position. To avoid this type of interference the system ensures that methods called by other methods in the same process use different ports. It does this by maintaining a pool of ports which were opened when the system started. As many ports are opened as are likely to be

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need d in th cours of pr cessing. Before a method can access a fil it calls a system p rt all cation routine in the same way that it would have called the "open" r utine. The p rt allocation routine finds a pre-opened port satisfying the type of access asked for and marks it as in use by this particular method. When the method completes, the system automatically marks the port free again (i.e., de-allocates it). If another method is called within the first method, then a different port will be allocated to it. The mapping of pooled ports to real ports is performed by the sam substitute file operation subroutines that are responsible for detecting transaction protection conflicts outlined above.

# B. Distributed Object-based Data Processing Applied to a Televised Auction System

Application of a distributed object-based data processing system to a televised auction system in accordance with a preferr dembodiment requires the following objects. The object's nod so (locations) and class are set forth adjacent to each of the objects in the table.

OBJECT NAME	NODES	CLASS
Obdic	Host	Obdic
Sessions	Host	Session
Users	Host	User
Auctions	Host	Auction
Currencies	Host	Currency
Sale	Host	Sale

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OBJECT NAME	NODES	CLASS
Sale Number	Host and Controller's W rkstation	Root
Bidlevel	Host, Currency Board Operator's Workstation, and Auctioneer's Display	Bidlevel
Currency board	Host, Television Display, and Saleroom Display	Root
Currentlot	Host, Controller's Workstation, Currency Board Operator's	Root
	Workstation, Auctioneer's Display, Television Display, and Saleroom Display	
Bidding Flag	Host, Controller's Workstation, and Currency Board Operator's Workstation	Root
Ping-pong flag	Host and Controller's Workstation	Root
Number of bidders	Host, Controller's Workstation, and Auctioneer's Display	Root
Leading bidder	Host, Controller's Workstation, and Auctioneer's Display	Root

OBJECT NAME	NODES	CLASS
Last accepted bidder	H st, Controller's Workstation, Auctioneer's Display, Television Display, and Saleroom Display	Root
Biddisplay	Host and `Controller's Workstation	Biddisplay
Logintab	Host	Logintab
Nextlots	Host and Controller's Workstation	Root
Bids	Host	Bid

Exemplary objects as set forth above are defined as follows:

"Obdic" is the object dictionary and contains the nam s of all objects, their locations, and a routing table which tells the computer the location of all objects and how to find each of the locations. This object is a basic part of the distributed object-based data processing system.

"Sessions" is a list of log-in times for each user, and where that user is located (e.g., Helsinki) for each session.

"Users" is a list of those entitled to use the system. These include paid subscribers to the system and the staff f the concern operating the system.

"Auctions" is information about the items to be auctioned, such as a list of the lots being sold.

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"Currencies" is a list f th various curr nci s to be displayed and the exchange rates.

"Sale" is a s t of details f the current sal which is in progress. The information in this object may be obtained from the "auctions" object.

"Sale number" is the number assigned to the current sale.

A blank indicates that no auction is currently in progress.

"Bid level" is the amount of the current bid.

"Currency board" is a translation of the bid level int the various currencies contained in the "Currencies" object.

"Current lot" is the lot number and miscellaneous details
of the current lot being sold.

"Bidding flag" is an indication whether bidding is in progress (i.e., the auction is not between two lots).

"Ping-pong flag" is an indication whether a ping-pong auction procedure is in progress.

"Number of bidders" is the number of bidders which bid through the system in the current round. This is obtained by instructing the computer to accumulate the number of bid signals received in any given round.

"Leading bidder" is the identification of the user whos bid signal was first received by the computer (including th first bidder in a ping-pong) or was "promoted" from the bid display by the controller.

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"Last acc pted bidd r" is th identificati n of the bidder whos bid was acc pt d by th auctione r in the last bidding round.

"Bid display" is a list of the bidders in the order in which the bids were received by the computer. This list is preferably limited in size to the top ten bidders.

"Log in table" is a table of users which have logged in for the current session.

"Next lots" is a description of several subsequent lots to be auctioned.

incoming bid.

Exemplary classes for computer implementation of the televis d auction system described herein are as follows:

"Session" class contains methods for determining the period a user is logged onto the system. This includes steps for determining log-in and log-out of a user.

"User" class contains methods which perform operations n the user file or another file such as a group file. These operations are, for example, adding, deleting, or updating a user (subscriber), and getting a user file for other operations.

"Currency" class contains methods for adding, deleting and changing currencies and currency exchange rates and for effecting exchange rate calculations.

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"Aucti n" class c ntains methods for adding, del ting, and g tting both an auction and a lot and for indicating when the hammer has fallen (as when an aucti n is terminated by acceptance of a final bid) and the next lot is to be auctioned.

"Date" class determines the current date and perf rms date manipulations.

"Sale" class contains methods to start an auction by initializing relevant objects to a useful state. For exampl, the bid display object must be set to zero or blank values, and the number or bidders object must be set to zero. These methods may be invoked by pressing a "start auction" button on the Controller's Workstation.

"Bid display" class contains methods to control the bid display, which is the display of the top ten bidders. These methods also update the bid display by adding or canceling a bid.

"Bid level" class contains methods to set a new bid level. For example, when the auctioneer accepts a bid, h states the price of the bid, and the currency operator provides an input specifying that level and invoking a method to update the bid level object.

"Bid" class contains a bid method, which updates the bid display object and the number of bidders object, and a bid cancel method, which removes bidders from the bid display,

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replaces th canc led bidd r, and updat s th number of bidd rs.

"Login tabl " class maintains a list of th users, th log in table object, who have logged on the system.

In the preferred embodiment, the functions described above ar performed by a general purpose computer, such that sold under the tradename "stratus", and by personal computers, such as those using the Microsoft DOS operating system. The personal computers are programmed to advise the general purpose computer that they manag a copy of a particular object and are to be informed of changes to it (e.g., they will receive "set" requests informing them or a new value).

### BRIEF DESCRIPTION OF THE DRAWINGS

Figure 1 is a schematic diagram of the preferred hardware f r implementation of an interactive, televised auction.

Figure 2 is an illustration of the auctioneer's display.

Figure 3 is an illustration of the controller's display.

Figure 4 is an illustration of a currency controller's display.

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## DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

illustrates a combination f computers communication elem nts which may be used to c nduct a tel vis d auction, or other interactive event. A host computer 2 is programmed with the distributed object-based transaction system described above and which has been tailored for a televised auction. The host computer 2 receives input from a subscriber by receiving signals generated at a subscriber terminal 4. The system is capable of receiving input from a large number of subscribers, but a single subscriber has been shown in the figures for illustration. In the ordinary arrangement, the subscribers are located at large distances from each other and from the location of the live auction and are, thus, large distances from the host The preferred data connection between the subscriber terminal and the host computer is a telephone line, which is connected to a packet-switching network such as 6.

The subscriber also has a television 8 which receives broadcast signals by way of a satellite network 10, or oth r broadcast system. The screen of the subscriber television 10 is preferably divided into four areas. The first area 12 contains a number (the "paddle" number) identifying the bidder whose bid was accepted in the last round and the location of that bidder. This is preferably a display of the "last accepted bidder" object. Ar a 14 of the subscriber's television screen contains the amount of th last accepted bid in the selected currencies. This may be a

display of "currency board" object. Ar a 16 contains a video display of the article being auctioned, and area 18 c ntains a video display of the auction er.

The areas 12, 14, 16, and 18 are generated by a television mixer 20 which combines signals from the host computer which produce the displays of the currency board and last accepted bidder objects with video signals from one or more television cameras (not shown) in the auction room directed at the article being auctioned and the auctioneer. The display of objects generated by the host computer may be assembled by a television display generator 22 which supplies signals to the television mixer 20.

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Several other display devices are provided for use by personnel associated with the auction. These are the auctioneer's display 24, the display on the currency converter workstation 26, and the display on the controllers's workstation 28. Each of this is connected to the host computer for supply by the host with signals representing the selected objects required by this articles. The preferred connection between the host and the workstations is by way of an X25 switch 30.

The auctioneer's display only receives signals from which it generates a display. An example of the auctioneer's display is shown in figure 2 wherein the number of the lot being sold is displayed at 32, the last accepted bid level at 34, the last accepted bidder at 36, and the number of bidders at 38. This display may be on a video screen located near the auctioneer such

that it is easily seen and is preferably projected on a s mitransparent scr n locat d such that the aucti neer can vi w simultan ously both the bidders in the sal room and the display.

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The currency converter workstation 26 and the operator's workstation 28 are preferably personal computers capable of providing input to the system, including the host computer 2, relating to their functions in the conduct of the auction. In n embodiment, two personal computers are supplied with programs such that either may serve as the currency converter workstation or the operator's workstation. Alternatively, these workstations may be provided with programs exclusive to the selected function. The screens are preferably touch sensitive whereby touching a selected display feature invokes an appropriate method of the distribut d object-based transaction system.

Figure 3 illustrates the display associated with the controller's workstation 28. This display contains the last accepted bidder at 40, the leading bidder at 42, the number of bidders at 44, the top ten bidders at 46, the number of the current sale at 48, and the lot number at 50. Touching one of the butter at 46 causes that bidder to be promoted to the leading bidder at 42, and touching the leading bidder display at 42 causes that bidder is bid to be accepted. Acceptance of a bid identifies that bidder as the "last accepted bidder".

A "ping-pong" bat 52 indicates that a ping-pong procedur is being c nducted by the auction er, and a button 54 permits th

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operat r to rem v the ping-p ng symbol wh n th ping-pong is t rminat d. A "hamm r fallen" button 56 allows th operator to indicate that th auction er has compl ted the sale of a lot, and touching this button causes the host computer to invoke the appropriate method to record such and to make any necessary updates of relevant objects.

operator's workstation. This display shows the last accepted bid level at 58 and a series (in this example, ten) of pre-set increments above the last accepted bid at 60. Buttons 62 adjacent each of the increment indicators allow the increments themselves to be adjusted within the preset increment in case the auction r calls for a bid within the pre-set increments. A button 70 allows the currency board operator to set initial values. Touching this causes a keypad display to be superimposed on the display of figur 4 whereby the operator may key in the initial values. Box 72 displays the increment between the values shown in boxes 60, which in the illustration is £100. The current lot is displayed at 74.

Buttons 76 and 78 are provided for the case where the auctioneer calls for a price wholly outside the scale of th display. Button 76 causes the entire display to be shifted down by a preset amount, while button 78 causes the display to be shifted upward.

Button 80 allows the operator to exit from the display.

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As the auction proceeds, the auction er states the price of the bid based inter alia up n the number of bidders. The currency operator enters this price by t uching the button 60 having that price. The increment and number of prices displayed are designed to ensure that in the majority of cases, one of the areas 60 will show the price called for by the auctioneer. If the correct price is not displayed already, however, the currency operator uses the buttons 62, 76 or 78 to adjust the display until the correct price is displayed in one of the boxes 60. That price is then selected by the operator's touching that box, and this invokes methods to update the "last accepted bid" object with that value at all nodes.

The subscriber's terminal 4 allows the subscriber to interact with the auction being conducted in the saleroom and viewed on television 8. The terminal 4 may be of the type manufactured by Verifone and provides a slot for receiving an identification card (not shown) which has been provided to the subscriber by the operator of the system after appropriate credit investigations, or the like. When a subscriber wishes to participate in an auction, the card is "wiped" through the slot, and the terminal 4 reads the subscriber's information from the card. The subscriber is prompted by messages on display screen 66 to enter an identification number by way of key pad 68. The identification number is verified by an appropriate method in the host computer after which the subscriber is permitted to participate in the auction.

The subscrib r's terminal includ s butt ns or other input devices f ractivation by the subscriber t allow the subscriber to signal the host comput r that h wishes t bid r cancel a bid. For example, the subscriber's terminal has a "bid" button which signals the host computer that the subscriber has bid on the article being auctioned.

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An auction utilizing the above described methods and devic s would proceed as follows.

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The controller would sign on at the controller's workstation, and the currency board operator would sign on at the currency board operator's workstation. The host computer then performs initializing methods which prepare the system to conduct the auction or perform system maintenance functions by updating any of the files, such as "user", "auctions", or the like.

A user begins by wiping the card issued by the auction operator in the slot on the subscriber's terminal and entering the assigned password. The terminal generates the "pin\_login" method in the "sessions" class, and this method passes the user identification from the card and the password which has ben entered by the user to the host for verification.

The televised auction is begun by the controller's pressing the "start auction" button on the controller's workstation when th auction in the saleroom is also begun. This activates the "b gin sale" method which may be found in the "sale" class. The "next lot" method is called to set the "current lot" to "l", the number

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of bids to zero, and the ping-pong flag to z ro. The bid display is also initialized and the bidding flag is set to one to allow bids to be received.

When a remote bidder presses the "bid" button on the terminal 4, the host computer is requested to perform the "bid" method n the "bid" object which is this case is the identification of th bidder. The host also increases the "number of bidders" by one, sets the leading bidder by providing the identification of the bid to be received first, and updates the bid display. This process is followed for each subsequent bid.

When a bid is accepted by the auctioneer, the "leading bidder" is set and the identification of the bidder whose bid was accept d is moved to the "accepted bidder" location on the display such as at 42 in the display of figure 3. This is accomplished by th controller by his touching the proper one of the buttons 46, if th bidder is remote, or the button 43, if the accepted bidder is in the saleroom. The currency operator selects the proper button to display the "last accepted bid" level.

When the auctioneer begins a ping-pong auction, the "ping-pong" flag is set by the controller activating a "ping-pong" button and identifying the participants of the ping-pong. This causes the ping-pong symbol to be displayed on the various displays and permits only the participants of the ping-pong to be listed on the display as the last accepted bidder or the leading bidder. The identifications of the there bidders are placed in the bid display,

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but none can be a "l ading bidder" without the auctione r first terminating the ping-pong.

When th auctioner signals that the sale of the lot is completed, the "hammer fallen" method is invoked by the controller activating a "hammer fallen" button. This sets the bidding flag to zero, which indicates that no bids will be accepted by the heat computer.

The steps followed by the distributed object-based transaction system are as follows. Activation of an input device, such as by touching a touch sensitive screen causes a "request" to be generated. That request comprises a method and an object and is in essence a statement to the computer to perform the stated method on the stated object. The computer first looks at the object table, such as one contained in the "obdic" object described above with respect to the interactive auction system. This table allows the computer to identify the devices on which the object resides, such as the host computer and one of th workstations. As noted above, it is a feature of the distributed object-based transaction system that the objects may reside on one or more separate devices. The computer determines the best route to the various locations of the object from the table and sends the instruction to all appropriate nodes where the object resides. method is then performed on the object at all of the nodes on which the object resides. A reply is then sent if the selected method or riginal r qu st call d for a r ply.

It will be appreciated that a unique transaction system has ben d scrib d. Modifications within the sc pe of the appended claims will be apparent to those of skill in th art.

#### W Claim:

1. An interactiv system c mprising:

h st c mputer means for p rf rming data operations;

a plurality of workstation terminal means connected to said host computer for displaying data from said host computer and for providing input to said host computer,

television network means for transmitting video information to a plurality of subscriber video terminals; and

a plurality of subscriber data terminal means for supplying data to said host computer.

- z. An interactive network according to claim 1 further comprising mixing means for supplying said data from said host computer to said television network for combination with said video information.
- 3. An interactive system according to claim 2 wherein said television network comprises television signal broadcast means for supplying said video information to said subscriber video terminals and said subscriber terminal means comprises telephone transmission means for supplying said data from said subscriber terminal means to said host computer.
- 4. An interactive system according to claim 3 wherein said video information produces an image of an item being sold on each of said subscriber video terminals and said data contains information about the price of said item.

- 5. An interactive system acc rding t claim 4 wher in at l ast one of said workstation terminal means displays said information about the price of said item.
- 6. An interactive system according to claim 5 wherein said information about the price of said item comprises the last bid for said item which has been accepted in an auction.
- 7. A method for conducting a transaction comprising providing a host computer with data regarding an item, providing an input terminal to at least one subscriber for transmitting signals to said host computer,

providing a video terminal for displaying an image of said item to said subscriber,

wherein said signals indicate transaction information with respect to said item.

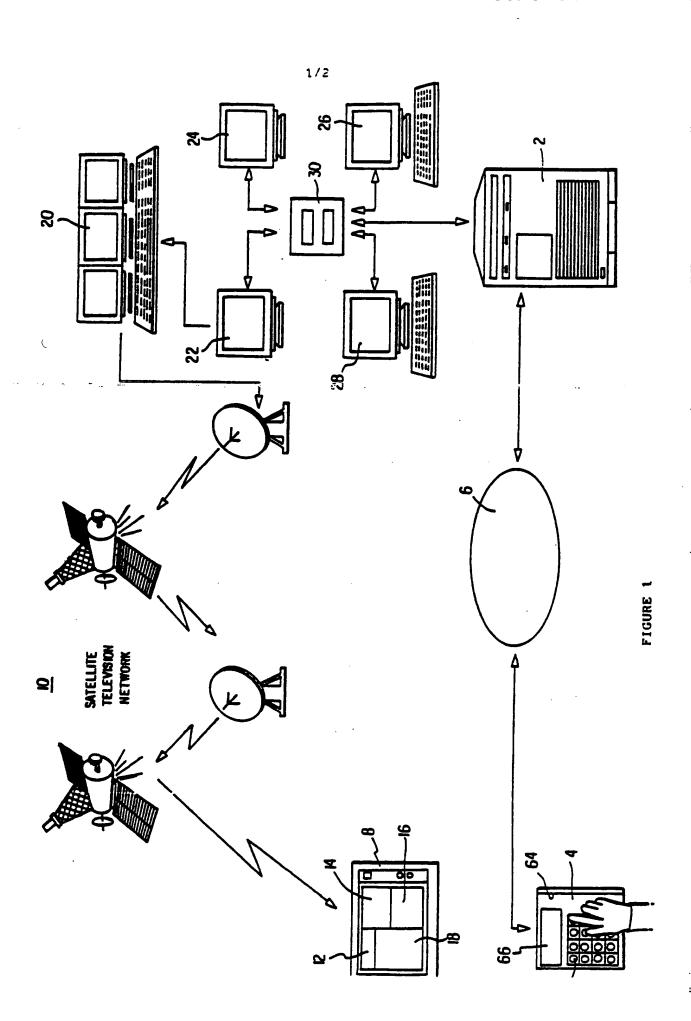
- 8. A transaction system comprising a host computer for performing data operations, a subscriber terminal for transmitting data from a subscriber to said host computer, video means for transmitting an image of an item to said subscriber, and a workstation terminal for displaying data from said host computer and for transmitting data to said host computer.
- 9. A method for processing data comprising providing a plurality of computing means,

defining a plurality of objects by allocating memory space in said computing means for each of said objects and by naming each f said objects,

pr viding at least one m thod for performing an operation with respect t said obj cts,

wherein said m mory spac f r at 1 ast on of said objects is allocated in a plurality of said computing means and said method includes the step of updating the value of said object at all memory spaces assigned to said object upon completion of said method.

10. A method according to claim 9 wherein said objects are related to an auction, and said at least one method comprises a plurality of methods relating to an auction.



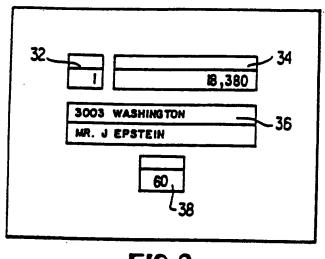
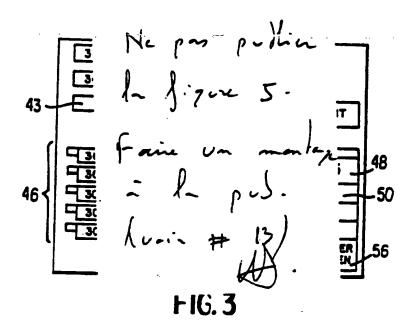


FIG. 2



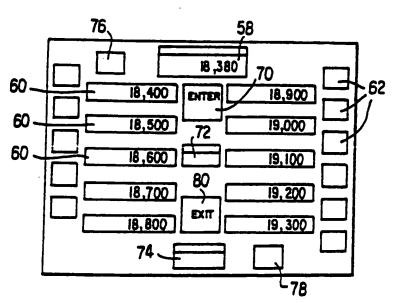


FIG. 4

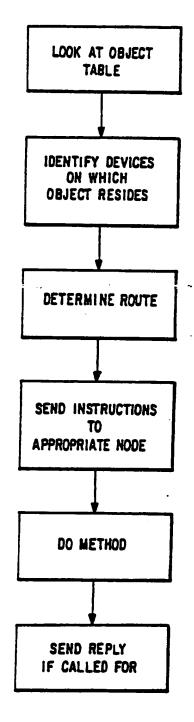


FIG. 5

Not to be taken in consideration for the international procedure

**Interactional Application No** L CLASSIFICATION OF SUBJECT MATTER (If several classification symbols apply, indicate all)6 According to International Patent Classification (IPC) or to both National Classification and IPC Int.C1. 5 HO4N7/173 II. FIELDS SEARCHED Minimum Documentation Searched? Classification System Classification Symbols Int.Cl. 5 H04N ; **H040** Decementation Searched other than Minimum Documentation to the Extent that such Documents are included in the Fields Searched III. DOCUMENTS CONSIDERED TO BE RELEVANT Citation of Document, 11 with indication, where appropriate, of the relevant passages 12 Relevant to Claim No.13 EP,A,O 342 295 (KATZ) 23 November 1989 1,7,8,9 - see the whole document 2-6.10 **FUNKSCHAU** 1,7,8,9 vol. 8, no. 19, September 1989, MUNCHEN pages 46 - 51; V.LUGGER: 'NEUES KONZEPT VERBINDET TV UND TELEFON' 2-5 see the whole document 1-5,7-9 16TH INTERNATIONAL TV SYMPOSIUM 16-22 JUNE 1989 MONTREUX, SWITZERLAND pages 554 - 559; S.ROBERTS ET AL.: 'VIDEOTEX AND INTERACTIVE SERVICES ON TREE&BRANCH CABLE TV NETWORKS' see page 557, line 6 - page 558, line 37 "I" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the \* Special categories of cited documents: 10 "A" document defining the general state of the art which is not considered to be of particular relevance "I" exciter document but published on or after the international "I" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to invelve an inventive step document which may threw doubts on priority citim(s) or which is cited to establish the publication date of another citation or other special reason (as specified) "Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such docu-ments, such combination being obvious to a person skilled "O" document referring to an eral discinsure, use, exhibition or other means document published prior to the interactional filling date but inter than the priority date claimed "A" document member of the same pittest family IV. CHRITIFICATION Date of Mailing of this International Search Report Date of the Actual Completion of the International Search 27. JS. 92 06 MAY 1992 Signature of Authorized Officer International Searching Authority GREVE M.P. Mac EUR PEAN PATENT OFFICE

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This assex lists the patent family members relating to the patent documents cited in the above-mentioned international search report. The members are as contained in the European Patent Office EDP file on The European Patent Office is in no way liable for those particulars which are merely given for the purpose of information. 06/05/92

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